

Model categories in Lean

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Source code

- ▶ <https://github.com/rwbarton/lean-model-categories>
Branch top-dev (under construction!)
- ▶ <https://github.com/rwbarton/lean-homotopy-theory>

What is a model category?

- ▶ (Quillen, *Homotopical Algebra*, 1967)
A *model category* is a category M equipped with three classes of morphisms
 - ▶ weak equivalences (W)
 - ▶ cofibrations (C)
 - ▶ fibrations (F)satisfying axioms **M1**, **M2**, **M3**, **M4**, **M5**.
- ▶ A model category is one kind of presentation of an underlying “homotopy theory” (or $(\infty, 1)$ -category) $L_W M$, which depends only on the weak equivalences W .
- ▶ The extra structure on M allows us to do certain calculations in $L_W M$ within the framework of ordinary categories.

Examples

- ▶ Classical homotopy theory of spaces
 - ▶ Topological spaces ($\text{Top}_{\text{Quillen}}$)

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 - ▶ $\text{Ch}(\text{RMod})_{\text{proj}}$, $\text{Ch}(\text{RMod})_{\text{inj}}$

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- ▶ Stable homotopy theory
- ▶ Equivariant homotopy theory
- ▶ Higher category theory, e.g., $(\infty, 1)$ -categories
- ▶ \mathbb{A}^1 -homotopy theory
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468 questions on MathOverflow with tag [model-categories]

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(23 questions containing the word “perfectoid”)

Project goals

- ▶ Formalize the definition of a model category (easy)
- ▶ Construct some of the central examples (hard)
 - ▶ Current focus: topological spaces ($\text{Top}_{\text{Quillen}}$)
 - ▶ Future aims: combinatorial model categories, simplicial sets

Definition: preliminaries

```
universe u
variables (M : Type (u+1)) [large_category M]

def morphism_class : Type* :=
   $\prod \{ \{ a \ b : M \} \}, (a \rightarrow b) \rightarrow \text{Prop}$ 

instance : has_subset (morphism_class M) :=
  { subset :=  $\lambda$  I J,
     $\forall \{ \{ a \ b : M \} \} \{ \{ f : a \rightarrow b \} \}, I \ f \rightarrow J \ f \}$  }

instance : has_inter (morphism_class M) :=
  { inter :=  $\lambda$  I J,  $\lambda$  a b f, I f  $\wedge$  J f }
```

Definition: weak equivalences

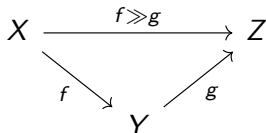
```
structure is_weq (W : morphism_class M) : Prop :=
  (weq_of_iso :  $\forall \{x y\} (i : \text{iso } x y), W i.\text{hom}$ )
  (weq_comp :  $\forall \{x y z\} \{f : x \rightarrow y\} \{g : y \rightarrow z\},$   

     $W f \rightarrow W g \rightarrow W (f \gg g)$ )
  (weq_cancel_left :  $\forall \{x y z\} \{f : x \rightarrow y\} \{g : y \rightarrow z\},$   

     $W f \rightarrow W (f \gg g) \rightarrow W g$ )
  (weq_cancel_right :  $\forall \{x y z\} \{f : x \rightarrow y\} \{g : y \rightarrow z\},$   

     $W g \rightarrow W (f \gg g) \rightarrow W f$ )
```

The last three conditions together are the “two-out-of-three property”.



In $\text{Top}_{\text{Quillen}}$, $W =$ maps inducing isomorphisms on π_* .

Definition: lifting property

```
def lp {a b x y : M} (f : a → b) (g : x → y) : Prop :=  
  ∀ (h : a → x) (k : b → y), h ≫ g = f ≫ k →  
  ∃ (l : b → x), f ≫ l = h ∧ l ≫ g = k
```

```
def llp (R : morphism_class M) : morphism_class M :=  
  λ a b f, ∀ {{x y}} {{g : x → y}}, R g → lp f g
```

```
def rlp (L : morphism_class M) : morphism_class M :=  
  λ x y g, ∀ {{a b}} {{f : a → b}}, L f → lp f g
```

$$\text{lp}(f, g) \iff \begin{array}{ccc} A & \xrightarrow{h} & X \\ f \downarrow & \exists l \nearrow & \downarrow g \\ B & \xrightarrow{k} & Y \end{array}$$

Definition: weak factorization system

```
structure is_wfs (L R : morphism_class M) : Prop :=  
  (llp : L = llp R)  
  (rlp : R = rlp L)  
  (fact :  $\forall \{x y\} (g : x \rightarrow y),$   
     $\exists z (l : x \rightarrow z) (r : z \rightarrow y), L l \wedge R r \wedge l \gg r = g)$ 
```

Factorization axiom: Any $g : X \rightarrow Y$ can be written as a composition $X \xrightarrow{l} Z \xrightarrow{r} Y$ with $l \in L$ and $r \in R$.

Definition: model category

```
class model_category (M : Type (u+1)) [large_category M] :=
  (complete : has_limits M)
  (cocomplete : has_colimits M)
  (W C F : morphism_class M)
  (weq : is_weq W)
  (caf : is_wfs C (F ∩ W))
  (acf : is_wfs (C ∩ W) F)
```


Constructing model categories

Construction of a model category structure on a category M usually proceeds along the following lines.

- ▶ Write down a class of weak equivalences W on M and check that W satisfies the two-out-of-three property.
- ▶ Construct two weak factorization systems (C, AF) and (AC, F) on M . There is a standard procedure for this called the **small object argument**.
- ▶ Verify that $AC = C \cap W$ and $AF = F \cap W$. This step generally requires some specialized argument.

Here, we'll focus exclusively on the second step.

Constructing weak factorization systems

- ▶ Recall that a weak factorization system (L, R) is required to satisfy $L = \text{llp}(R)$ and $R = \text{rlp}(L)$.
- ▶ Easy lemma: $\text{rlp}(\text{llp}(\text{rlp}(I))) = \text{rlp}(I)$.
(Same as constructive proof of $\neg\neg\neg P \iff \neg P$.)
- ▶ Thus, we can choose any collection of morphisms I and set $L = \text{llp}(\text{rlp}(I))$ and $R = \text{rlp}(I)$, and then the conditions $L = \text{llp}(R)$ and $R = \text{rlp}(L)$ will be satisfied. We say that (L, R) is “cofibrantly generated” by I .

For $\text{Top}_{\text{Quillen}}$, we will take

- ▶ (C, AF) to be generated by $I = \{ \partial D^n \rightarrow D^n \mid n \geq 0 \}$
- ▶ (AC, F) to be generated by $J = \{ D^n \times \{0\} \rightarrow D^n \times [0, 1] \mid n \geq 0 \}$

Constructing weak factorization systems

However, we still need to verify the last condition for a weak factorization system, namely that every morphism g can be factored as a morphism of L followed by a morphism of R .

We'll specialize to this situation:

- ▶ $M = \text{Top}$
- ▶ $I = \{ \partial D^n \rightarrow D^n \mid n \geq 0 \}$
- ▶ $Y = *$ (so $g : X \rightarrow *$)

Need to express g as a composition of $j : X \rightarrow X'$ and $g' : X' \rightarrow *$ with $j \in \text{llp}(\text{rlp}(I))$ and $g' \in \text{rlp}(I)$.

The small object argument

- ▶ Suppose g already has the property that every map $\partial D^n \rightarrow X$ admits an extension to D^n .

$$\begin{array}{ccc} \partial D^n & \longrightarrow & X \\ \downarrow & \nearrow \exists & \downarrow \\ D^n & \longrightarrow & * \end{array}$$

Then $g \in \text{rlp}(I)$ and we can factor g as the identity followed by g .

- ▶ In general, this will not be the case. Idea: replace $g : X \rightarrow *$ with a new map $g' : X' \rightarrow *$ which is “closer” to having this property.

The small object argument

- Form X' by attaching a cell D^n along every possible attaching map $\partial D^n \rightarrow X$. This yields a factorization of $g : X \rightarrow *$ as

$$X \xrightarrow{j} X' \xrightarrow{g'} *$$

By construction, any map $\partial D^n \rightarrow X$ admits an extension landing in X' :

$$\begin{array}{ccccc} \partial D^n & \longrightarrow & X & \xrightarrow{j} & X' \\ \downarrow & & \searrow \exists & \searrow g & \downarrow g' \\ D^n & \xrightarrow{\quad} & & & * \end{array}$$

- General fact: $j \in \text{llp}(\text{rlp}(I)) = L$.
- Done?

The small object argument

- ▶ There will be new maps $\partial D^n \rightarrow X'$, so that $g' \notin \text{rlp}(I)$.
- ▶ Repeat! Set $X_{i+1} = X'_i$, obtaining a sequence

$$X = X_0 \xrightarrow{j_0} X_1 \xrightarrow{j_1} X_2 \xrightarrow{j_2} \dots$$

- ▶ Define $X_\omega = \text{colim}_i X_i$.
- ▶ General fact: the composition $j : X \rightarrow X_\omega$ still belongs to $\text{lfp}(\text{rlp}(I)) = L$.

The small object argument

- ▶ Fact (not general!): Any map $\partial D^n \rightarrow X_\omega$ factors through X_i for some $i < \omega$.
- ▶ That means any lifting problem for $g_\omega : X_\omega \rightarrow *$ factors through $g_i : X_i \rightarrow *$ for some i , and therefore we attached a solution to this lifting problem at stage $i + 1$.

$$\begin{array}{ccccccc} \partial D^n & \longrightarrow & X_i & \xrightarrow{j_i} & X_{i+1} & \longrightarrow & X_\omega \\ \downarrow & & \exists & \nearrow & \searrow & & \downarrow g_\omega \\ & & D^n & \xrightarrow{\quad} & & & * \\ & & & \nearrow g_i & & & \end{array}$$

- ▶ Therefore, $g_\omega \in \text{rlp}(I) = R$ and we're done.

The transfinite case

- ▶ In general, countably many steps may not be sufficient.
- ▶ We can extend this construction transfinitely:
 - ▶ At a successor stage, construct $X_{\alpha+1}$ from X_α as before.
 - ▶ At a limit stage, define X_β to be the colimit of the entire sequence $(X_\alpha)_{\alpha < \beta}$ defined so far.
- ▶ In broad generality, one can show that there exists an ordinal γ such that every map $A_i \rightarrow X_\gamma$ factors through X_α for some $\alpha < \gamma$. Then we can end the construction at stage γ .

Formalizing the small object argument

- ▶ We need to construct the transfinite sequence $X_0 \rightarrow X_1 \rightarrow \dots \rightarrow X_\omega \rightarrow X_{\omega+1} \dots \rightarrow X_\gamma$.
- ▶ Obvious idea: some kind of recursive construction.
- ▶ Choose a well-ordered type α with order type $\gamma + 1$.
- ▶

```
def well_founded.fix {α : Sort u} {C : α -> Sort v}
  {r : α -> α -> Prop} (hwf : well_founded r)
  (F : Π x, (Π y, r y x -> C y) -> C x)
  (x : α) : C x
```
- ▶ How to choose C ? Maybe take $C\ i$ to be the type of sequences $X_0 \rightarrow X_1 \rightarrow \dots \rightarrow X_i$?

Formalizing the small object argument

- ▶ At a limit stage β , we need to form the “union” $(X_\alpha)_{\alpha < \beta}$ of all the sequences defined earlier, so that we can define X_β to be its colimit. In order to make sense of this union, we must know that each sequence extends the earlier ones.
- ▶ However, the induction hypothesis $\prod y, y < x \rightarrow C y$ of `well_founded.fix` does not give us any information about how the previous values were defined.
- ▶ *If* we could perform this recursive construction, then we could use the defining equations for `well_founded.fix` to prove that each sequence extends the earlier ones...

Recursive constructions with `roption`

- ▶ Idea: Recursively construct a sequence of values, each of which is either the desired thing, or an “undefined” value.

```
structure roption ( $\alpha$  : Type u) : Type u :=  
  (dom : Prop)  
  (get : dom ->  $\alpha$ )
```

Classically, `dom` is either `true`, in which case `get` provides a value of type α , or `false`, in which case `get` provides nothing. So `roption α` is isomorphic to `option α` .

- ▶ After constructing these `roption` values recursively, prove by induction that all the `dom` fields are actually `true`.

General setup

```
parameters { $\alpha$  : Type u} {r :  $\alpha \rightarrow \alpha \rightarrow \text{Prop}$ }  
  (hwf : well_founded r)  
local infix `<` := r
```

```
parameters {C :  $\alpha \rightarrow \text{Type u}$ }  
parameters (q :  $\prod \{i\ j\}, C\ i \rightarrow C\ j \rightarrow \text{Prop}$ )  
local infix `~` := q
```

```
parameters  
  (F :  $\prod (a : \alpha) (f : \prod (i : \alpha), i < a \rightarrow C\ i)$   
    (hf :  $\forall i\ j\ ria\ rja, i < j \rightarrow f\ i\ ria \sim f\ j\ rja$ ),  
    {x : C a //  $\forall i (ria : i < a), f\ i\ ria \sim x$ })
```

- ▶ C is the major premise, i.e., what we want to construct at each step.
- ▶ q is a compatibility condition on the values of C at each pair of steps i, j with $i < j$.
- ▶ F describes how to perform each inductive step.

The recursive definition

```
def crec_core :  $\Pi (a : \alpha), \text{roption } (C a) :=$   
hwf.fix $  $\lambda a I,$   
{ dom :=  $\exists (t : \forall i (ria : i < a), (I i ria).dom),$   
   $\forall i j ria rja, i < j \rightarrow$   
   $(I i ria).get (t i ria) \sim (I j rja).get (t j rja),$   
  get :=  $\lambda h,$   
  F a ( $\lambda i ria, (I i ria).get (h.fst i ria) h.snd$ ) }
```

- ▶ The value at step a is defined if the values at all steps $i < a$ are defined and consistent.
- ▶ In that case, we can use the inductive step F to produce the new value.

The definition is actually total

```
lemma crec_lemma :  $\forall$  (a :  $\alpha$ ),  
   $\exists$  (t : (crec_core a).dom)  
    (t' :  $\forall$  i (ria : i < a), (crec_core i).dom)  
    (hq :  $\forall$  {{i j ria rja}} {{rij : i < j}},  
      (crec_core i).get (t' i ria)  $\sim$   
      (crec_core j).get (t' j rja)),  
    (crec_core a).get t =  
      (F a ( $\lambda$  m rma, (crec_core m).get (t' m rma)) hq).val  $\wedge$   
     $\forall$  i (ria : i < a),  
      (crec_core i).get (t' i ria)  $\sim$  (crec_core a).get t :=  
  ...
```

- ▶ Straightforward inductive argument.
- ▶ Key point: (t : (crec_core a).dom), so crec_core a has a value.

Extracting the value

```
def crec (a :  $\alpha$ ) : C a :=  
  (crec_core a).get (crec_lemma a).fst  
  
lemma crec_consistent {i j} (h : i < j) :  
  crec i ~ crec j :=  
  by apply (crec_lemma j).snd.snd.snd.right; exact h
```

Generic way to carry out an inductive construction while maintaining a consistency invariant.

Current status of $\text{Top}_{\text{Quillen}}$

```
axiom A : rlp serre_I = rlp serre_J  $\cap$  @is_weak_equivalence
```

```
def quillen_serre : model_category.{1 0} Top.{0} :=  
model_category.mk' W_is_weq serre_caf serre_acf A AC_sub_W
```


Thank you!